



Approximate Homomorphic Encryption

- Construction & Bootstrapping

Yongsoo Song, Seoul National Univ

ECC 2018, Osaka



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Microsoft Research, Redmond

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- **Related Works**

Advanced Cryptography

- Protecting Computation, not just data



Advanced Cryptography

- Protecting Computation, not just data
- Differential Privacy
- Zero-knowledge Proof
- Multiparty Computation
- Attribute Based Encryption
- ...

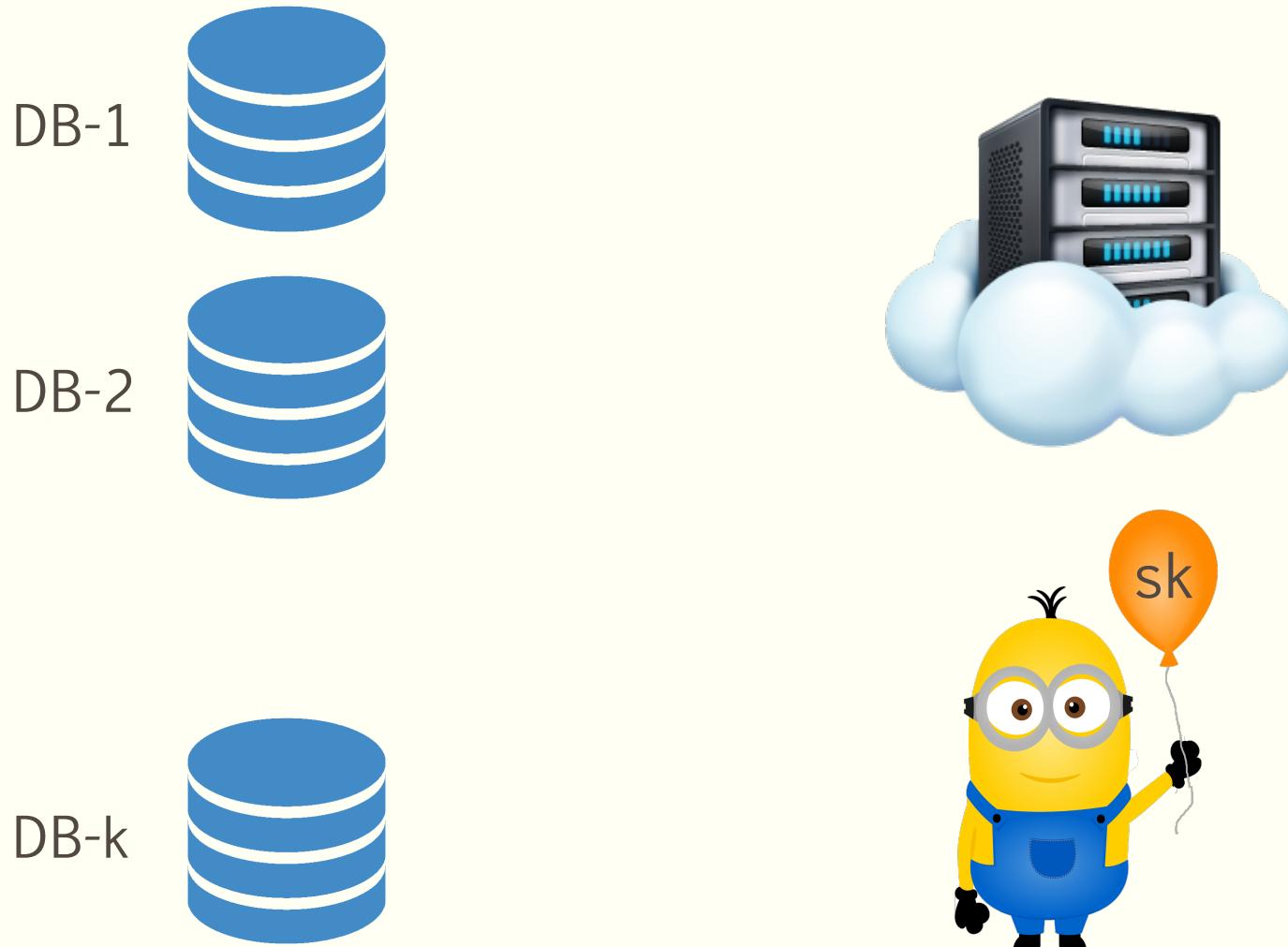


Advanced Cryptography

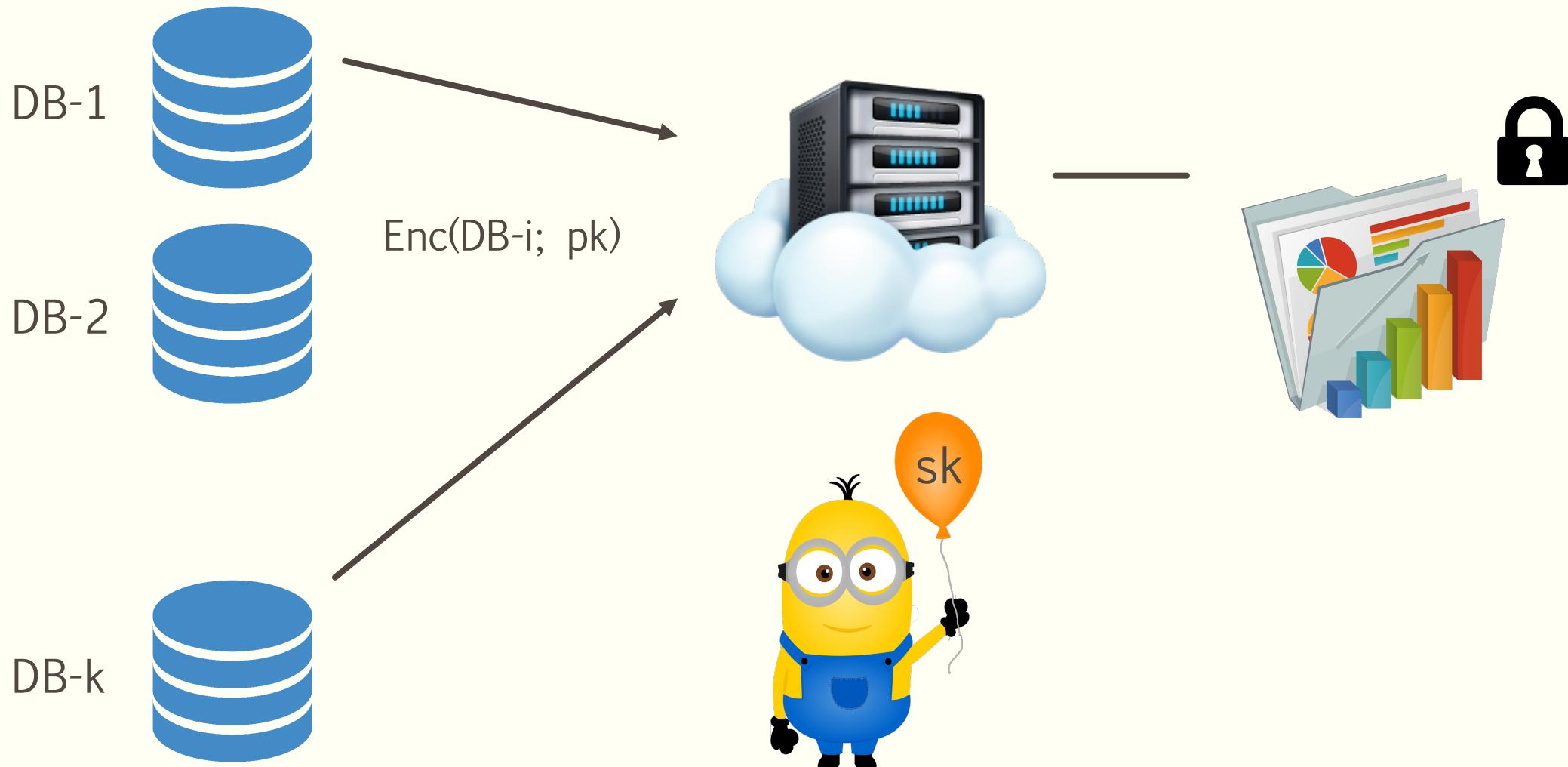
- Protecting Computation, not just data
- Differential Privacy
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- Attribute Based Encryption
- ...
- Homomorphic Encryption (2009~)



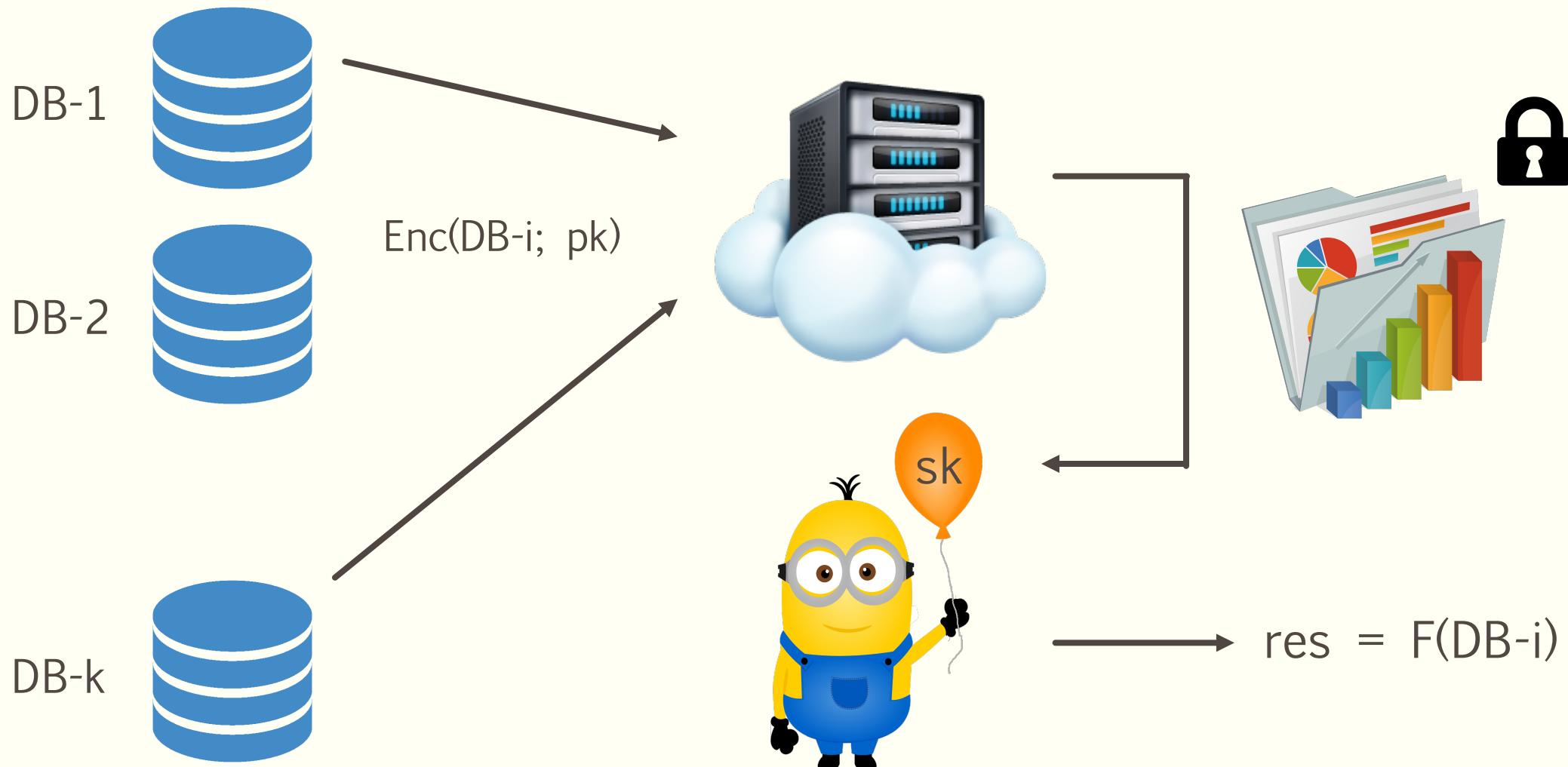
Homomorphic Encryption



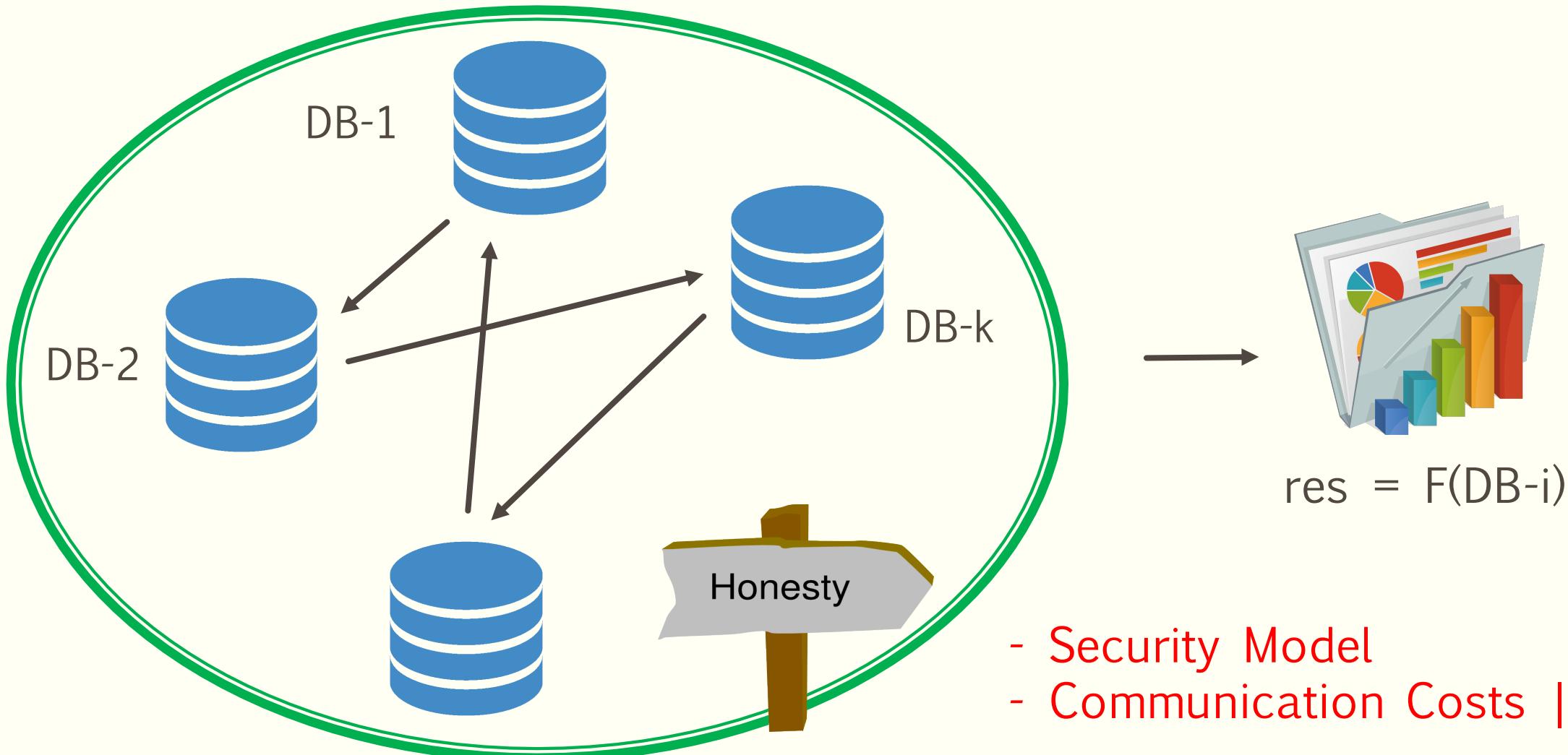
Homomorphic Encryption



Homomorphic Encryption



Multi-Party Computation



Comparison: HE vs MPC

	Homomorphic Encryption	Multi-Party Computation
Re-usability	One-time encryption No further interaction	Single-use encryption Interaction between parties each time
Model	Semi-honest Cloud + Trusted SK Owner	Semi-honest parties without collusion
Speed	Slow in computation (but can speed-up using SIMD)	Slow in communication (due to large circuit to be exchanged)

Summary of Progresses

- 2009-10: Plausibility
 - [GH11] A single bit operation takes 30 minutes
- 2011-12: Large Circuits
 - [GHS12b] 120 blocks of AES-128 (30K gates) in 36 hours



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- 2015-today: Usability
 - Various schemes with different advantages
 - Simpler and faster implementations
 - Real-world tasks: Big data analysis, Machine learning
 - Standardization meetings (2017~)
 - iDASH competitions (2014~)



4 Big Takeaways from Satya Nadella's Talk at Microsoft Build



By [JONATHAN VANIAN](#) May 7, 2018

[Microsoft](#) CEO Satya Nadella is trying to distinguish the business technology giant from its technology brethren by focusing on digital privacy.

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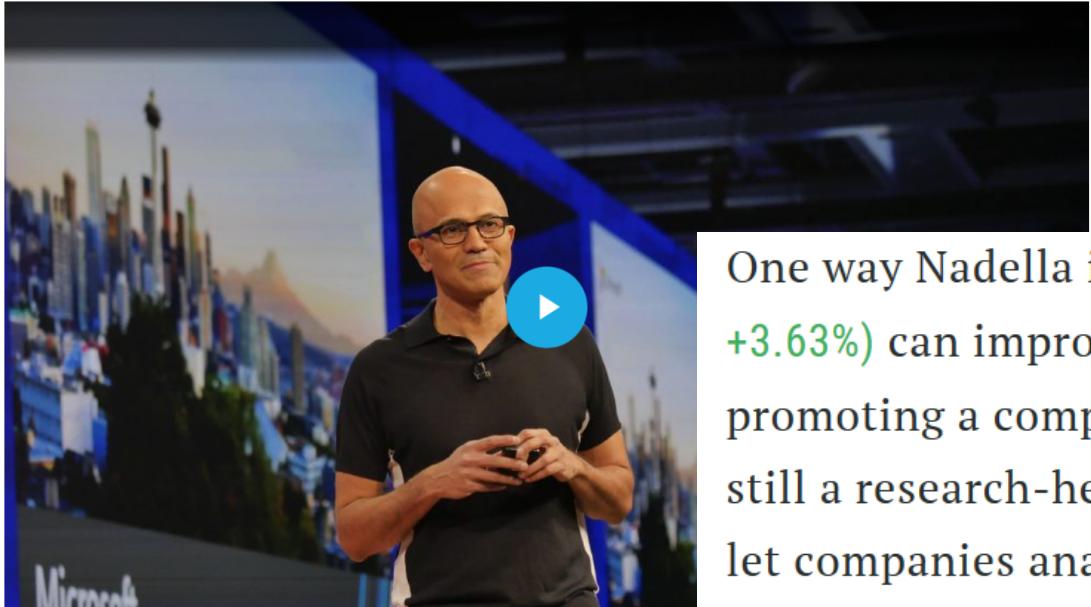
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Microsoft CEO Satya Nadella is trying to distinguish the giant from its technology brethren by focusing on digital

One way Nadella is attempting to convince businesses that Microsoft ([MSFT, +3.63%](#)) can improve its AI technology while protecting user data is by promoting a computing technique called homomorphic encryption. Although still a research-heavy technique, [homomorphic encryption](#) would presumably let companies analyze and crunch encrypted data without needing to unscramble that information.

Nadella is pitching the technique as a way for companies to [“learn, train on encrypted data.”](#) The executive didn’t explain how far along Microsoft is on advancing the encryption technique, but the fact that he mentioned the wonky terms shows that the company is touting user privacy as a selling point for its Azure cloud business.

Best Performing HE Schemes

Type	Classical HE	Fast Bootstrapping	Approximate Encryption
Scheme	[BGV12] BGV [Bra12, FV12] B/FV	[DM15] FHEW [CGGI16] TFHE	[CKKS17] HEAAN
Plaintext			
Operation			
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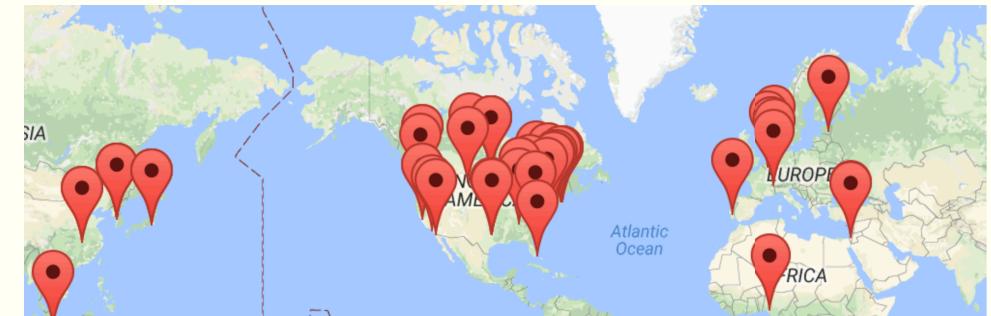
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Plaintext	Finite Field Packing	Binary string	Real/Complex numbers Packing
Operation	Addition, Multiplication	Look-up table & bootstrapping	Fixed-point Arithmetic
Library	HElib (IBM) SEAL (Microsoft Research) Palisade (Duality inc.)	TFHE (inpher, gemalto, etc.)	HEAAN (SNU)

iDASH Security & Privacy Workshop

- An interdisciplinary challenge on genomic privacy research
- Motivated by real world biomedical applications
- Participation of privacy technology experts (academia and industry)
- Developed practical yet rigorous solutions for privacy preserving genomic data sharing and analysis
- Reported in the media (e.g., Nature News, GenomeWeb)



nature international weekly journal of science

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Archive > Volume 519 > Issue 7544 > News > Article

NATURE | NEWS

Extreme cryptography paves way to personalized medicine

Encrypted analysis of data in the cloud would allow secure access to sensitive information.

Erika Check Hayden

23 March 2015

PDF | Rights & Permissions

Cloud processing of DNA sequence data promises to speed up discovery of disease-linked gene variants.

David Paul Morris/Bloomberg via Getty

genomeweb

Home | Business & Policy | Technology | Research | Clinical | Disease Areas

me = The Scan | To Keep It Safe and Sound

Keep It Safe and Sound

25 March 2015

One of the concerns about using genetic data along with medical records to personalize medicine is how to keep that personal information safe, if it still easily accessible for analysis. Cryptographers at a workshop hosted by the University of California, San Diego, **tested a homomorphic encryption method** that promises, reports *Nature News* Erika Check Hayden.

This method involves mathematically encrypting data on a local computer and then loading the encoded form to the cloud where it can be analyzed, Check Hayden notes. Encoded results are then sent back to a local computer, which unscrambles the data. Any data intercepted along the way would be encrypted.

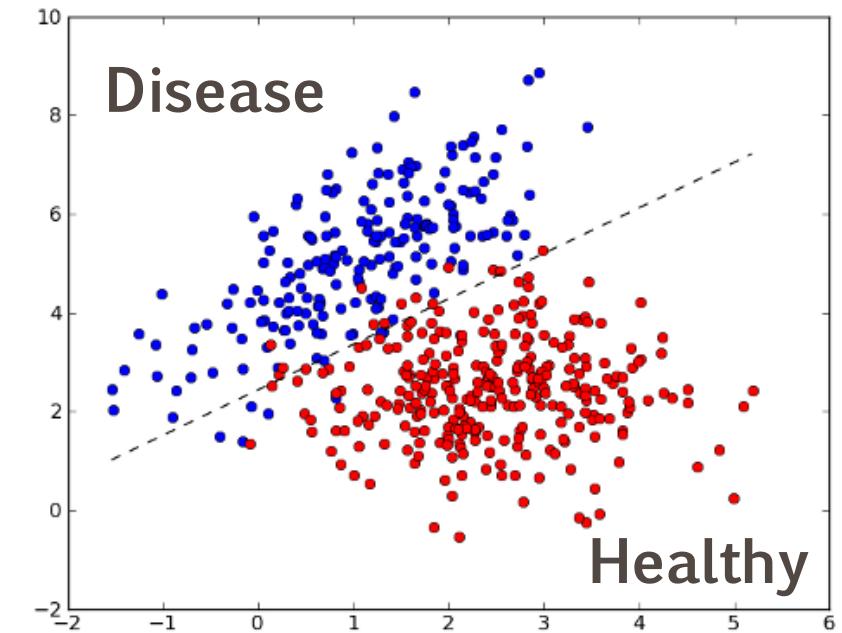
She notes that this idea dates back to 1978, but remained largely theoretical until 2009 when IBM Thomas J. Watson Research Center's Craig Gentry showed that computational analyses could be carried out on homomorphically encrypted data.

At the UCSD workshop, cryptographers showed that such an approach could analyze data from 400 people within about 10 minutes and pinpoint a variant associated with disease from among few hundred loci. Analysis of larger datasets of more base pairs wasn't always possible, Check Hayden says, and it could take lots of computer memory, time, or money.

While the workshop organizers find the approach promising, others say it might not provide enough protection for the data or allow researchers and clinicians to perform all the analyses they want. US National Center for Biotechnology Information's Steven Sherry, for instance, prefers restricting data access to a select few people who have agreed to follow certain regulations on how the data may be used.

iDASH 2017 – Logistic Regression Model Training

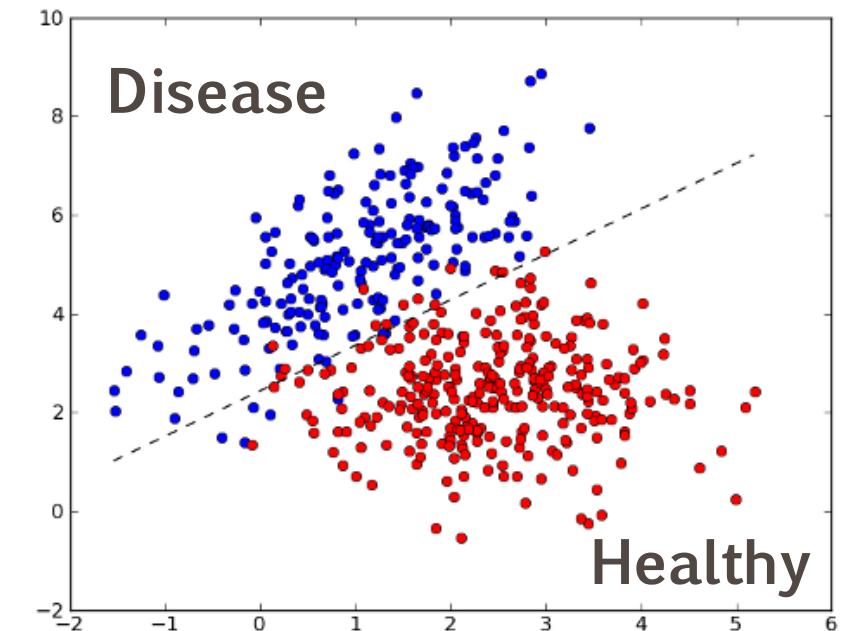
- A machine learning model to predict the disease
- 1500 records + 18 features for training



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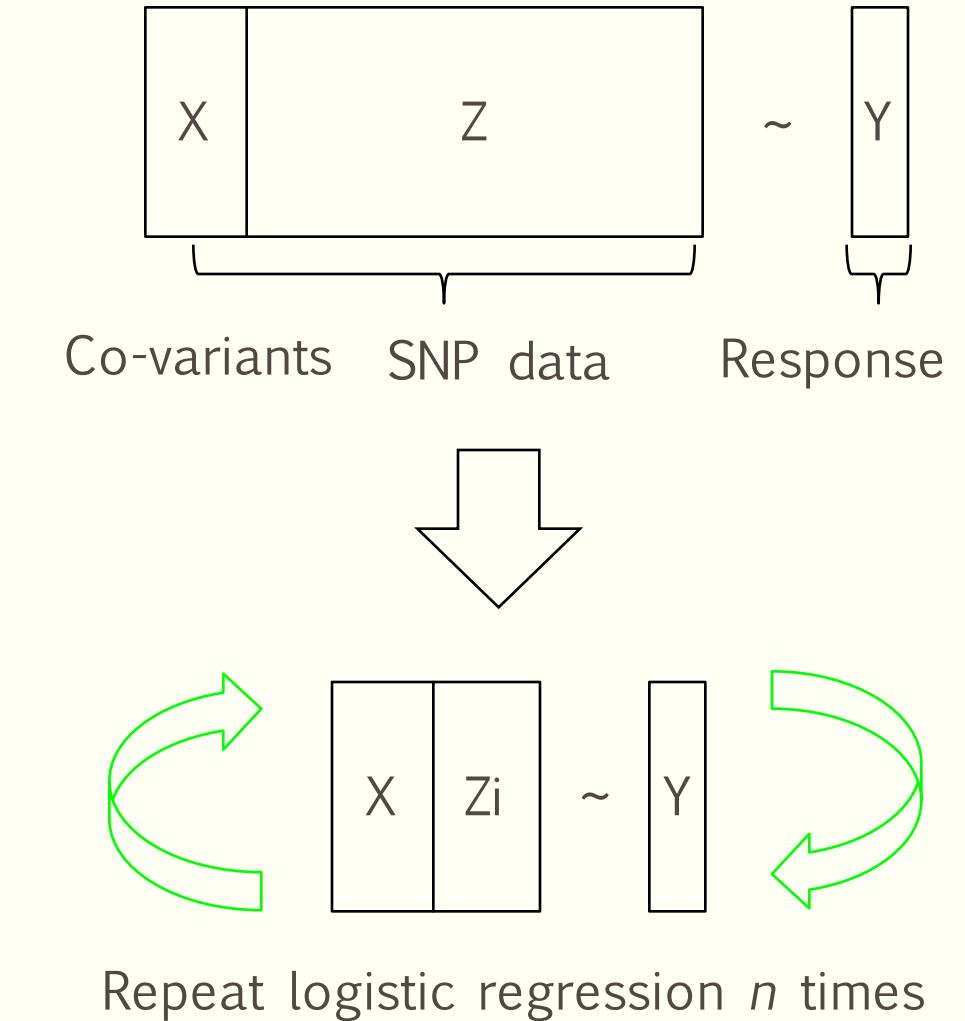
Teams	AUC	Secure learning		Overall time (mins)
		Time (mins)	Memory (MB)	
SNU	0.6934	10.250	2775.333	10.360
CEA LIST	0.6930	2206.057	238.255	2207.363
KU Leuven	0.6722	155.695	7266.727	160.912
EPFL	0.6584	15.089	1498.513	16.739
MSR	0.6574	385.021	26299.344	396.390
Waseda*	0.7154	2.077	7635.600	5.332
Saarland**	N/A	48.356	29752.527	57.344



* Interactive mechanism, no complete guarantee on 80-bit security at “analyst” side

iDASH 2018 – Semi-Parallel GWAS

- Compute Genome Wide Association Studies (GWAS)
- 3 Co-variants [age, height, weight] + 14,841 SNPs



iDASH 2018 – Semi-Parallel GWAS

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- 3 Co-variants [age, height, weight] + 14,841 SNPs

Team	Submission	Schemes	Time (mins)	Memory (MB)	Accuracy
A*FHE	A*FHE 1	HEAAN	922.48	3,777	0.999
	A*FHE 2		1,632.97	4,093	0.905
Chimera	Version 1	TFHE+HEAAN (Chimera)	201.73	10,375	0.993
	Version 2		215.95	15,166	0.35
Delft Blue	Delft Blue	HEAAN	1,844.82	10,814	0.969
UCSD	Log Reg	HEAAN	1.66	14,901	0.993
	Lin Reg	pkg: RNS HEAAN	0.42	3,387	0.989
Duality Inc	Log Reg	HEAAN	3.80	10,230	0.993
	Chi2 test	pkg: PALISADE	0.09	1,512	0.983
SNU	SNU 1	HEAAN	52.49	15,204	0.984
	SNU 2		52.37	15,177	0.988
IBM	IBM-Complex	HEAAN	23.35	8,651	0.911
	IBM- Real	pkg: HELib	52.65	15,613	0.526

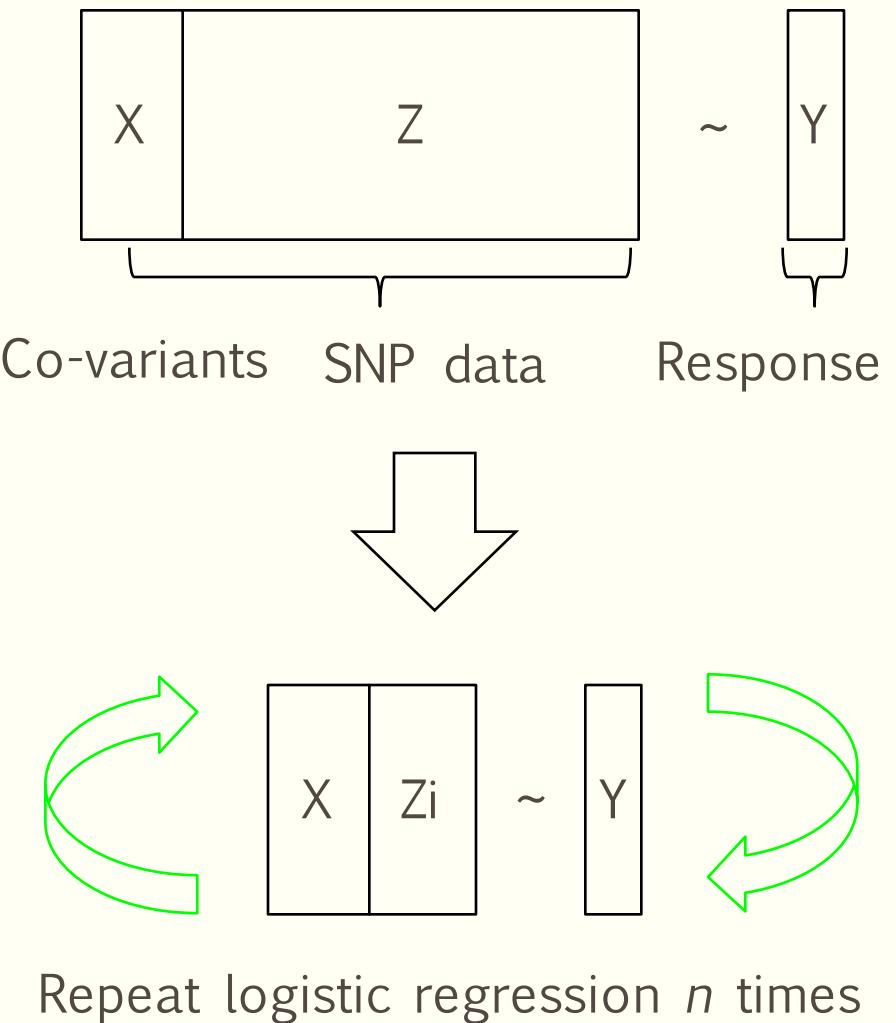


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Approximate Computation

- Numerical Representation

Encode m into an integer $m \approx px$ for a scaling factor p

$$\sqrt{2} \mapsto 1412 \approx \sqrt{2} \cdot 10^3$$

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- Fixed-Point Multiplication

Compute $m = m_1 m_2$ and extract its significant digits $m' \approx p^{-1} \cdot m$

$$1.234 \times 5.678 = (1234 \cdot 10^{-3}) \times (5678 \cdot 10^{-3}) = 7006652 \cdot 10^{-6} \xrightarrow{\text{red}} 7007 \cdot 10^{-3} = 7.007$$

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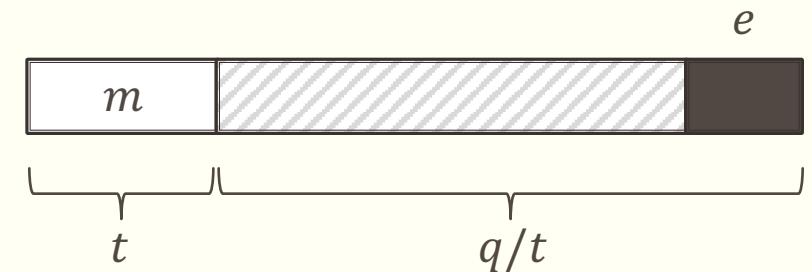
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- Previous (LWE-based) HE

$$\text{ct} = \text{Enc}_{\text{sk}}(m), \quad \langle \text{ct}, \text{sk} \rangle = \frac{q}{t}m + e \pmod{q}$$

Modulo t plaintext vs Rounding operation

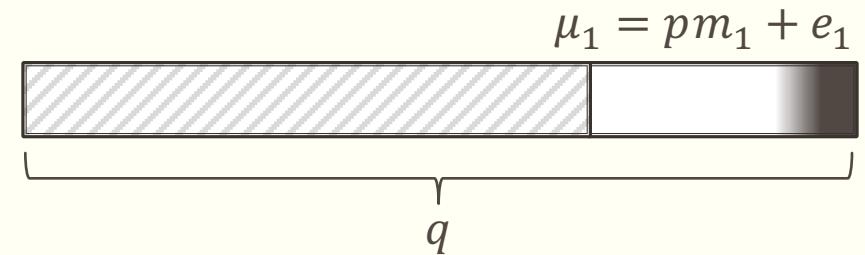


HEAAN

- A New Message Encoding

$$ct = \text{Enc}_{\text{sk}}(m), \langle ct, \text{sk} \rangle = pm + e \pmod{q}$$

Consider e as part of approximation error



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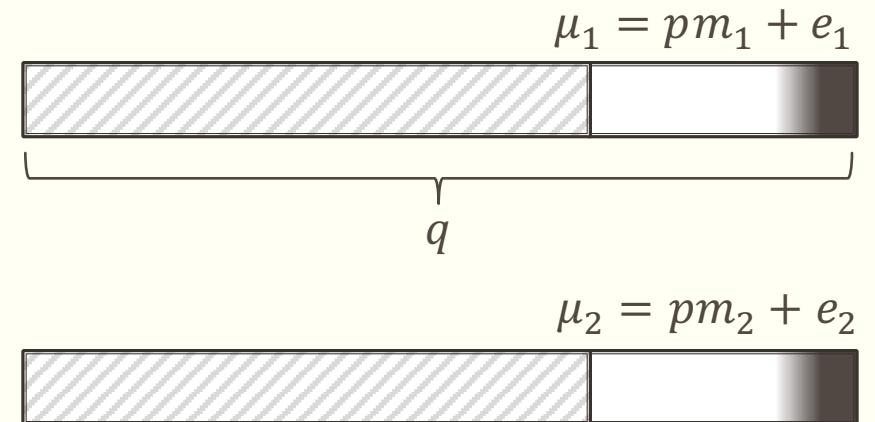
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- Homomorphic Operations

Input: $\mu_1 \approx pm_1, \mu_2 \approx pm_2$

Addition: $\mu_1 + \mu_2 \approx p \cdot (m_1 + m_2)$



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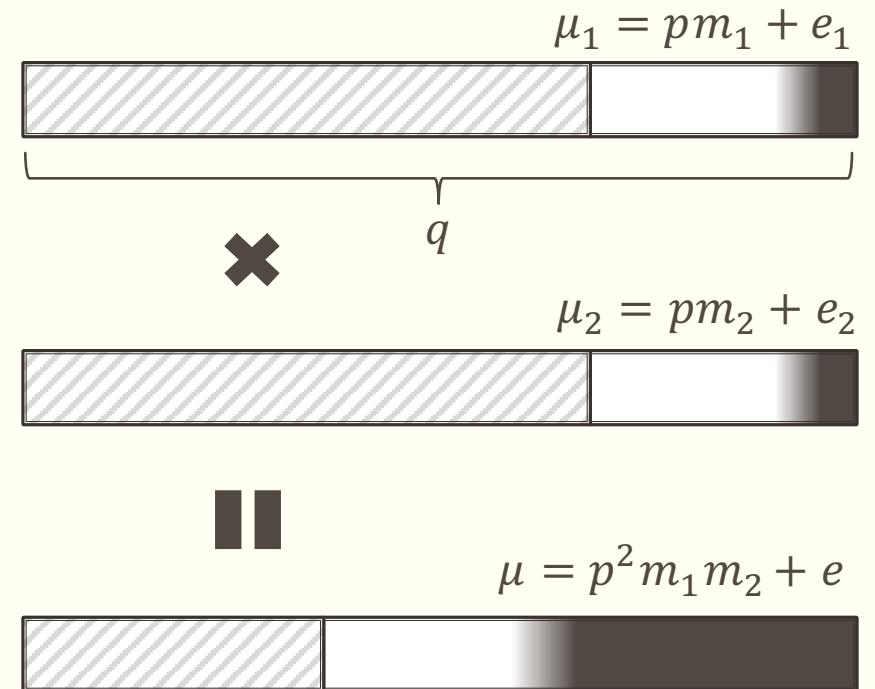
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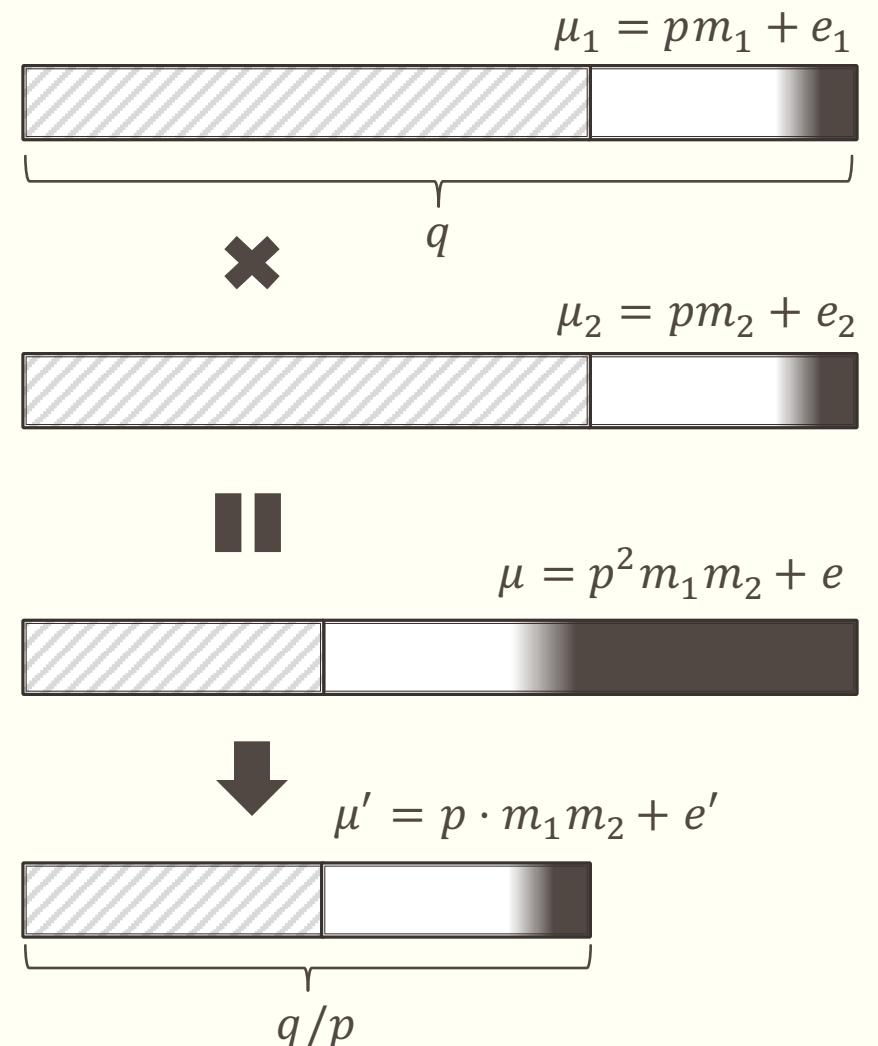
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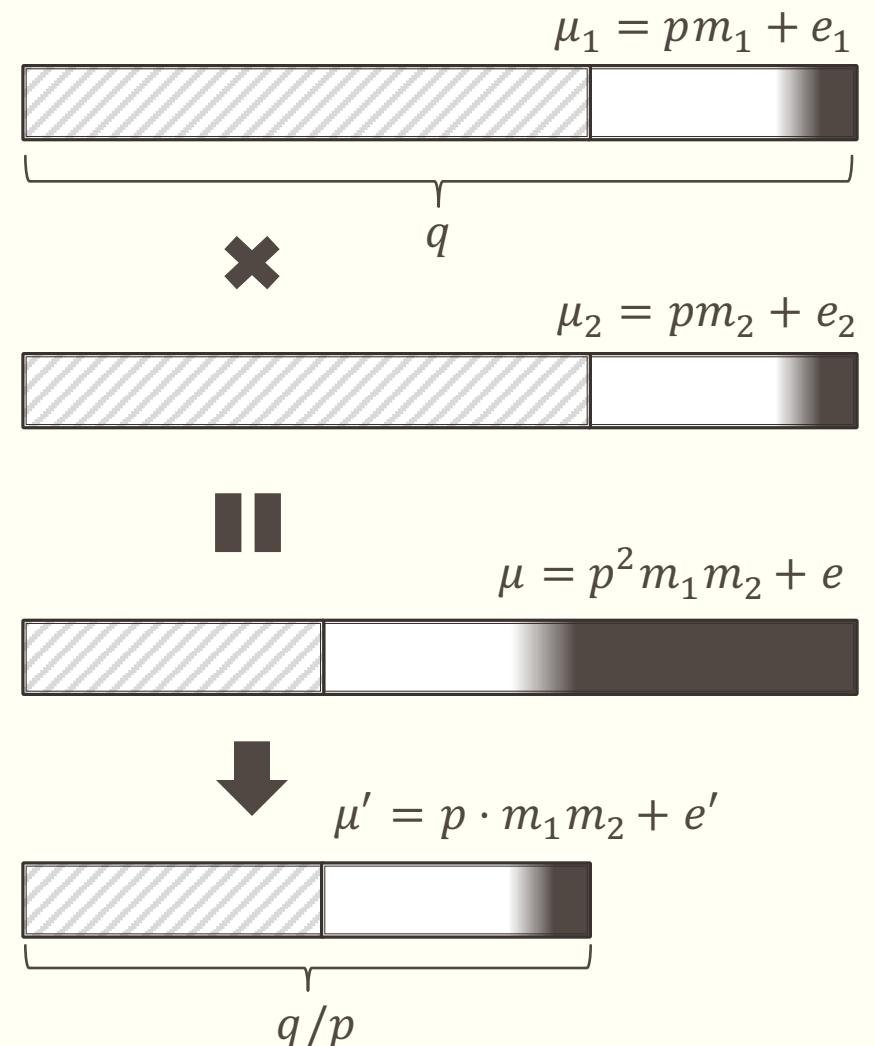
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- Support for the (approximate) fixed-point arithmetic !

- **Leveled HE** : $q = p^L$



Packed Ciphertext

- Construction over the ring $R = \mathbb{Z}[X]/(X^n + 1)$ and $R_q = R \pmod{q}$

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 - A single ciphertext can encrypt a vector of plaintext values $z = (z_1, z_2, \dots, z_\ell)$
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- RLWE-based HEAAN
 - A ciphertext can encrypt a polynomial $m(X) \in R$
 - Observation: $X^n + 1 = (X - \zeta_1)(X - \zeta_1^{-1})(X - \zeta_2)(X - \zeta_2^{-1}) \dots (X - \zeta_{n/2})(X - \zeta_{n/2}^{-1})$

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 - Decoding/Encoding function

$$R = \mathbb{Z}[X]/(X^n + 1) \subseteq \mathbb{R}[X]/(X^n + 1) \rightarrow \mathbb{C}^{n/2}$$

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- Permutation of plaintext slots

$$\text{Rotate: } \text{Enc}(z_1, z_2, \dots, z_{n/2}) \mapsto \text{Enc}(z_2, \dots, z_{n/2}, z_1)$$

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Bootstrapping of HEAAN

- Bootstrapping
 - Ciphertexts of a leveled HE have a limited lifespan
 - Refresh a ciphertext $ct = \text{Enc}_{\text{sk}}(m)$ by **evaluating the decryption circuit homomorphically**

$$\text{Dec}_{\text{sk}}(ct) = m \Leftrightarrow F(\text{sk}) = m \text{ where } F(*) = \text{Dec}_*(ct)$$

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- Bootstrapping key $\text{BK} = \text{Enc}_{\text{sk}}(\text{sk})$

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- HEAAN

- Homomorphic operations introduce errors

$$F(\text{BK}) = F(\text{Enc}_{\text{sk}}(\text{sk})) = \text{Enc}_{\text{sk}}(F(\text{sk}) + e) = \text{Enc}_{\text{sk}}(m + e)$$

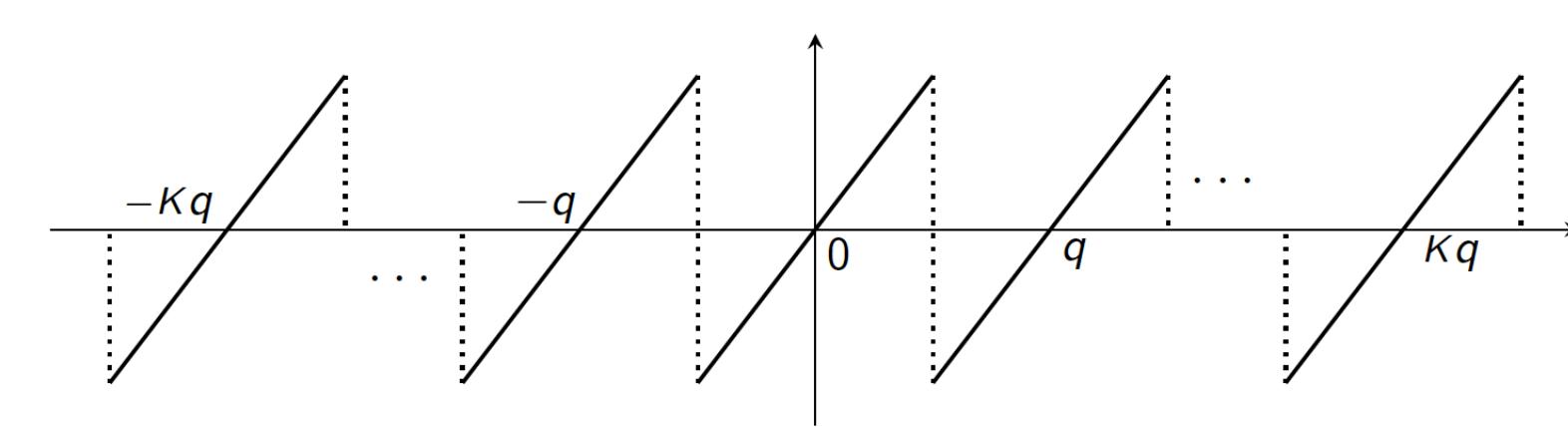
- It is ok to have an additional error
 - How to evaluate the decryption circuit efficiently?

$$\text{Dec}_{\text{sk}}(ct) = \langle ct, \text{sk} \rangle \pmod{q}$$

Approximate Decryption

$$\begin{aligned} \text{Dec}_{\text{sk}}(\text{ct}) &\mapsto t = \langle \text{ct}, \text{sk} \rangle \xrightarrow{\text{red}} [t]_q = \mu, \\ &\quad t = qI + \mu \text{ for some } |I| < K \end{aligned}$$

- Naïve solution: polynomial interpolation on $[-Kq, Kq]$
- Huge depth, complexity & inaccurate result

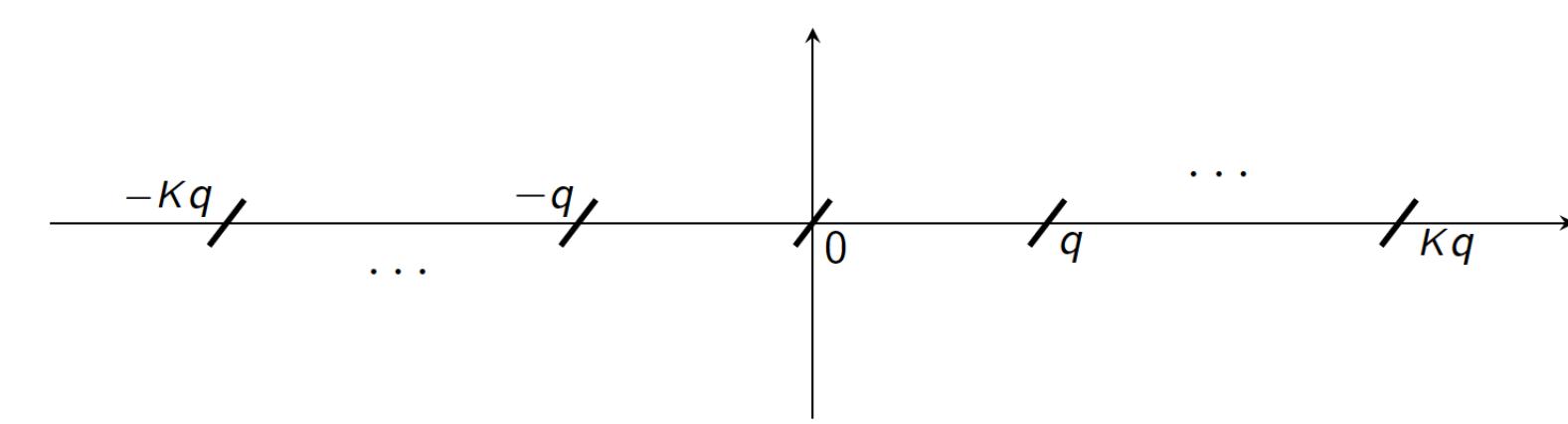


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- Idea 1: Restriction of domain $|\mu| \ll q$

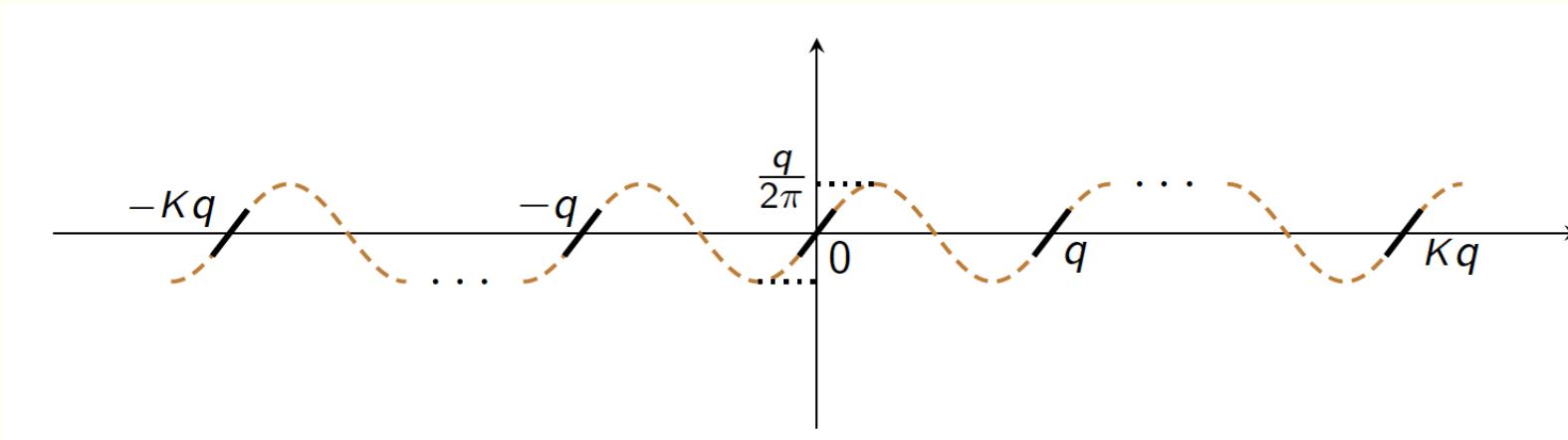


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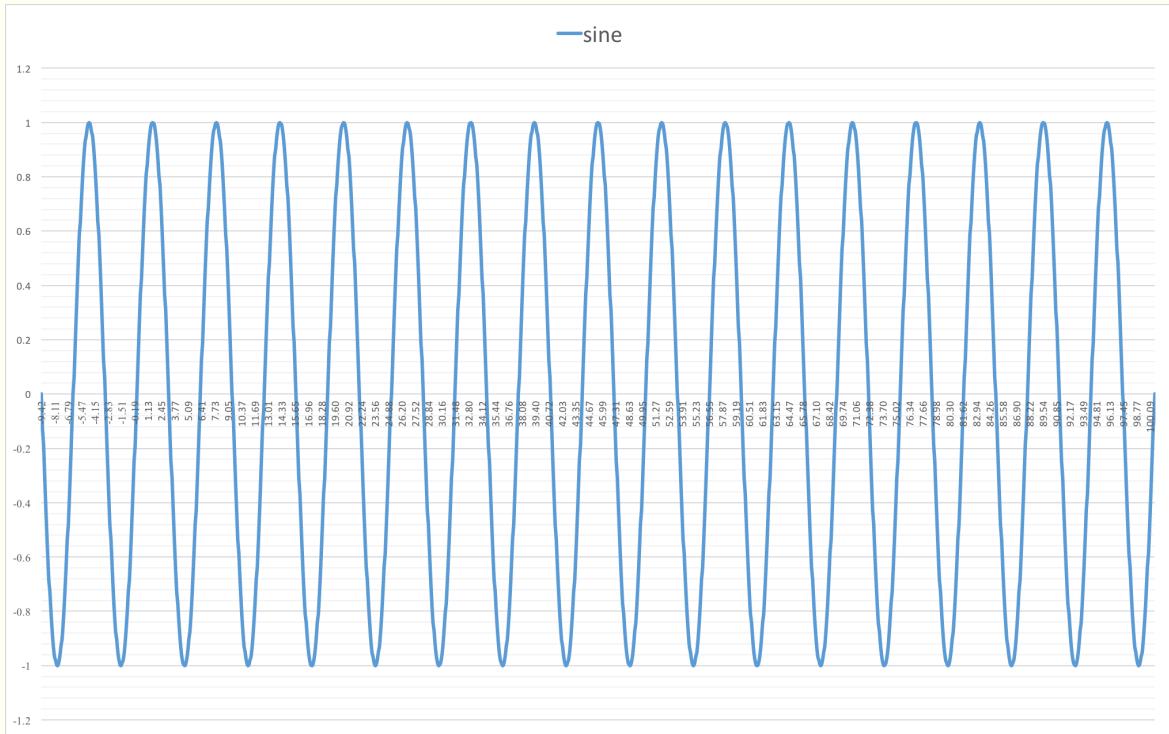
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- Idea 1: Restriction of domain $|\mu| \ll q$
- Idea 2: Sine approximation $\mu \approx \frac{q}{2\pi} \sin \theta$ for $\theta = \frac{2\pi}{q} t$

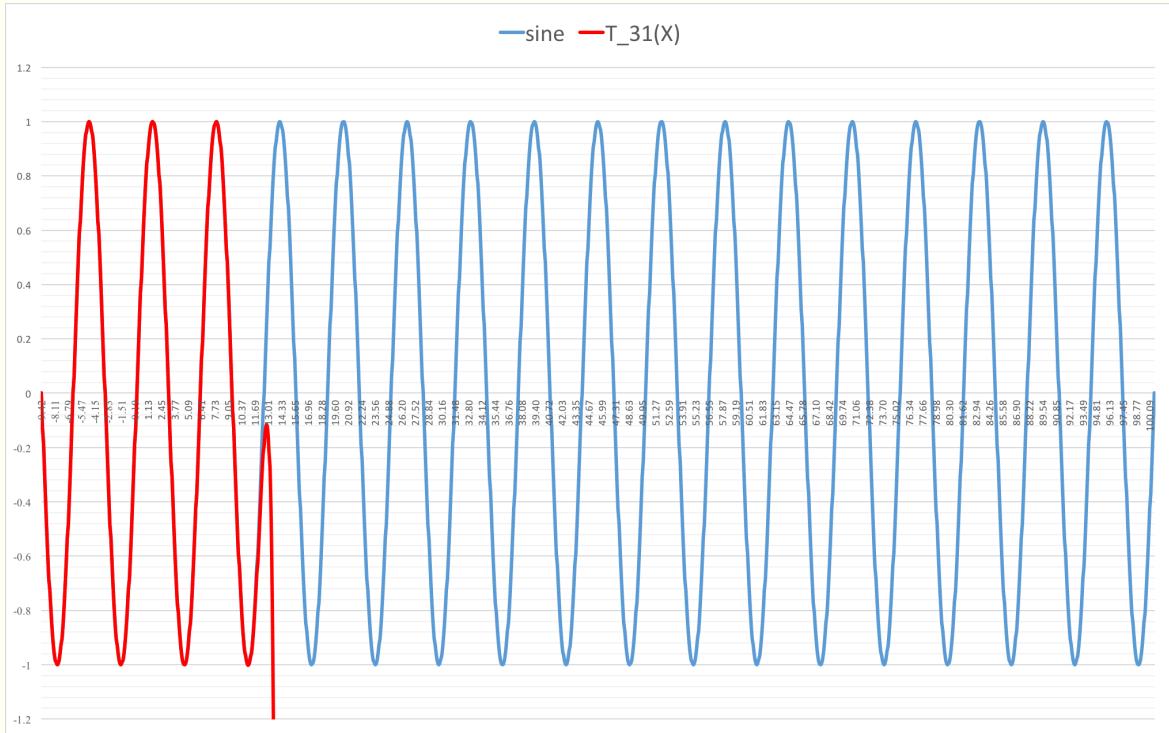


Sine Evaluation



Sine Evaluation

- Direct Taylor approximation
 - huge depth & complexity, low precision

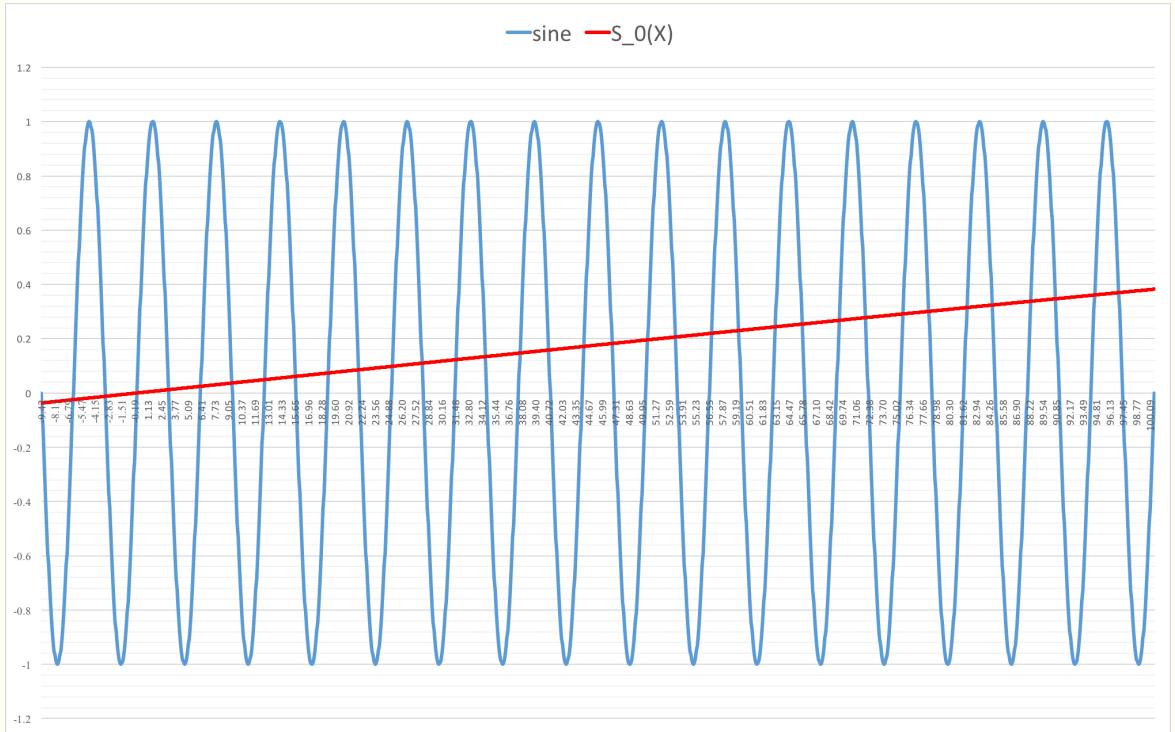


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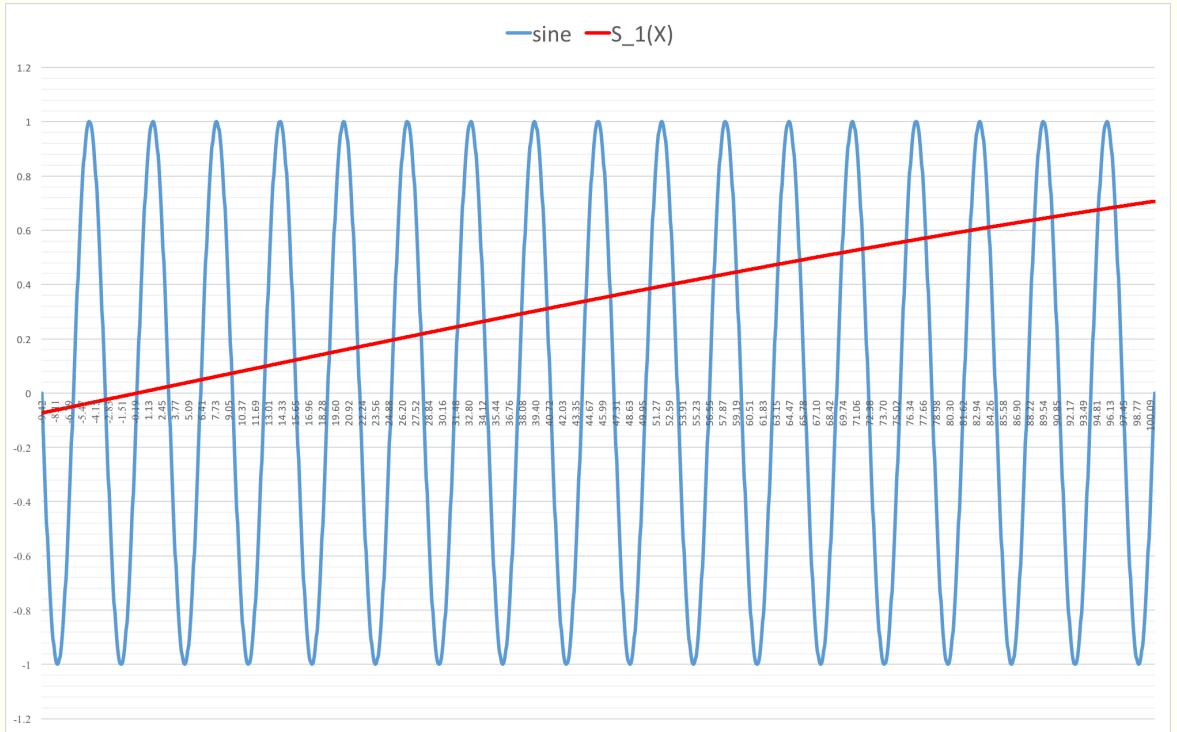
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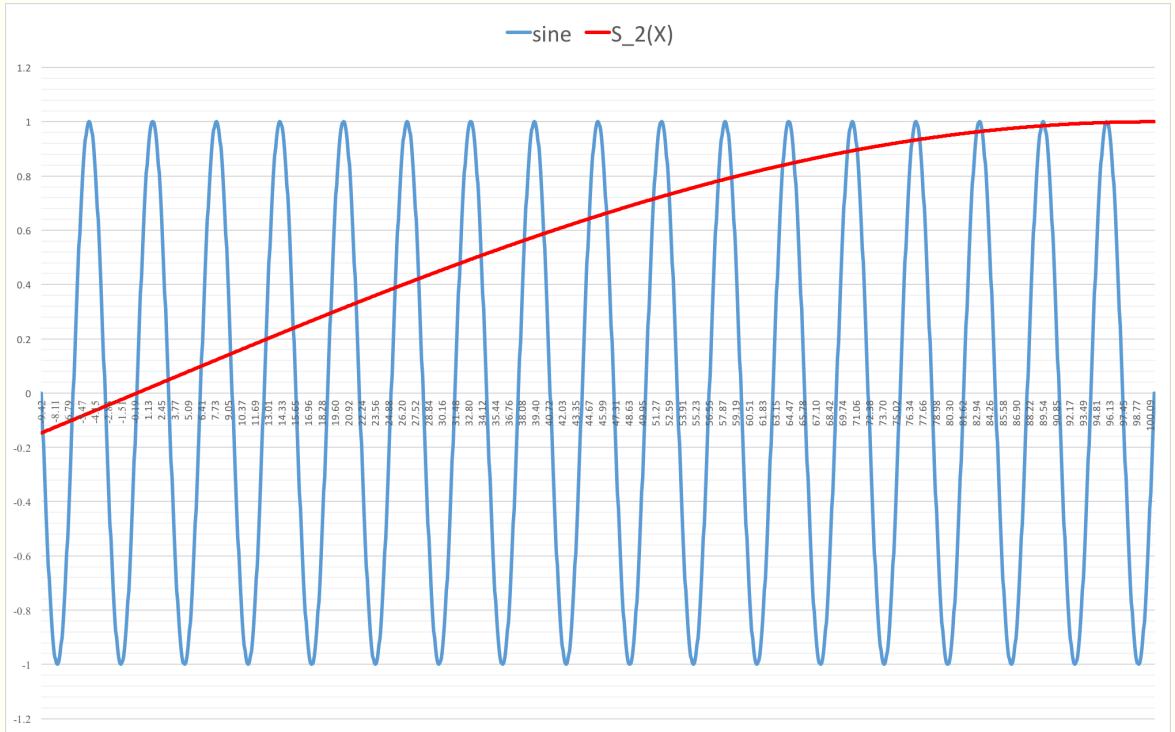
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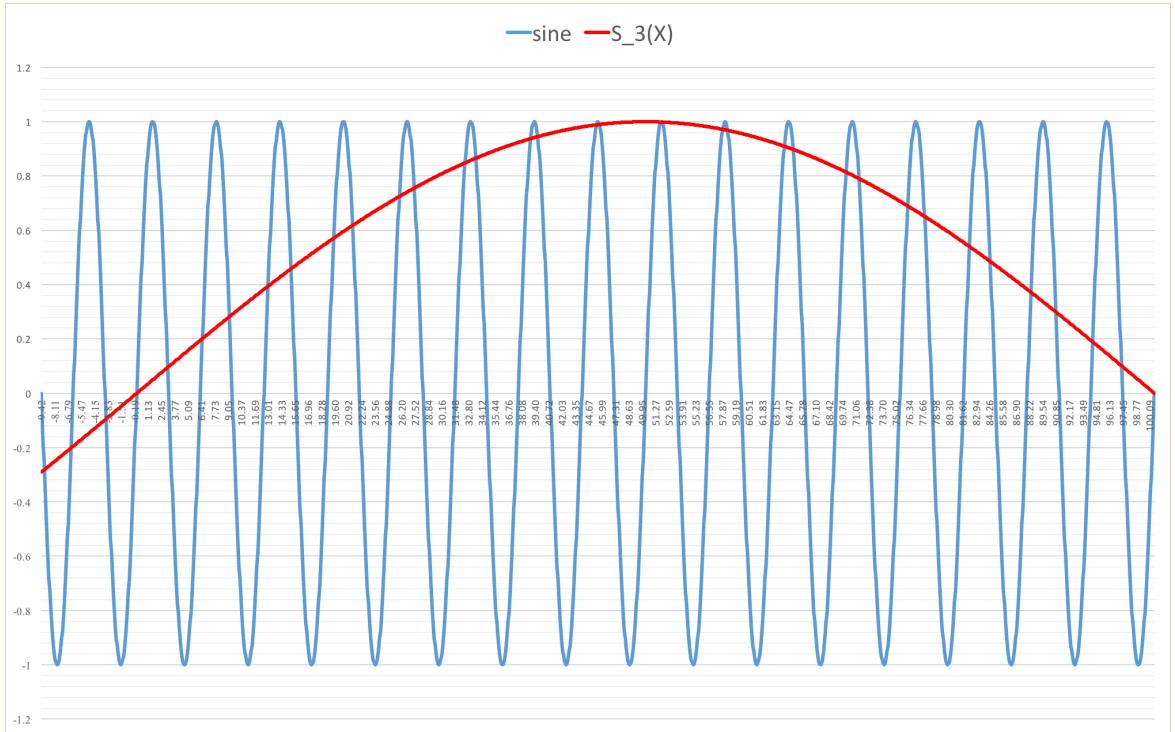
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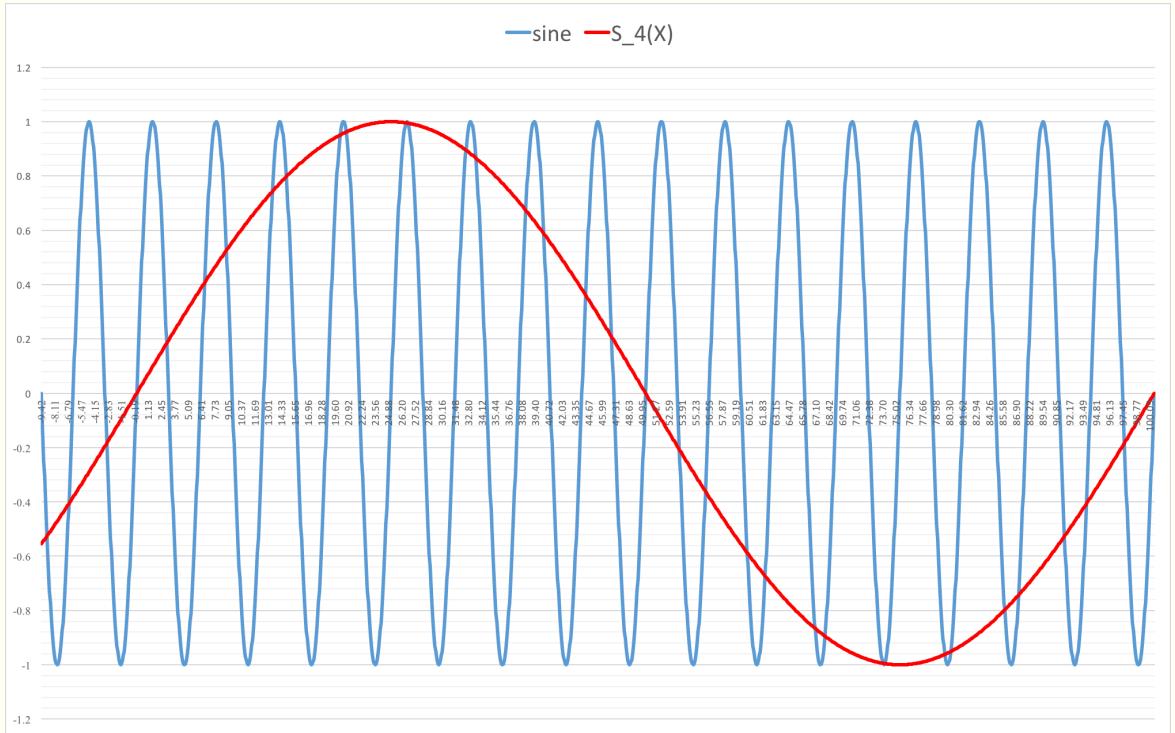
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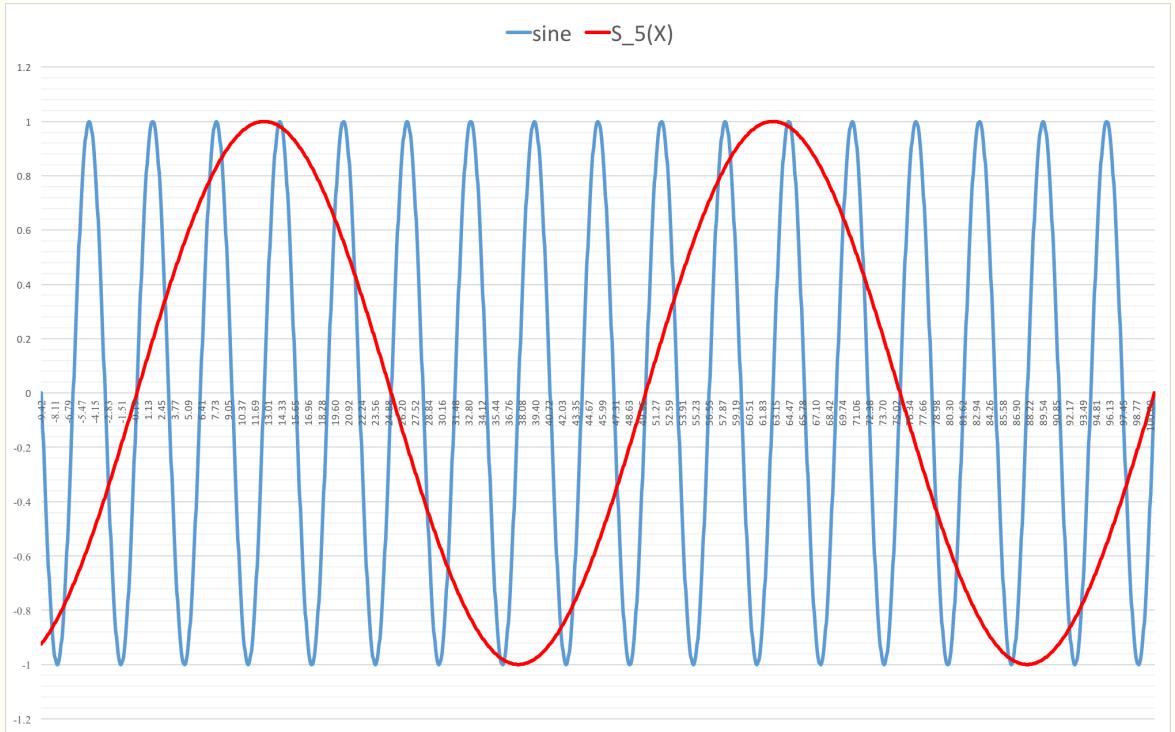
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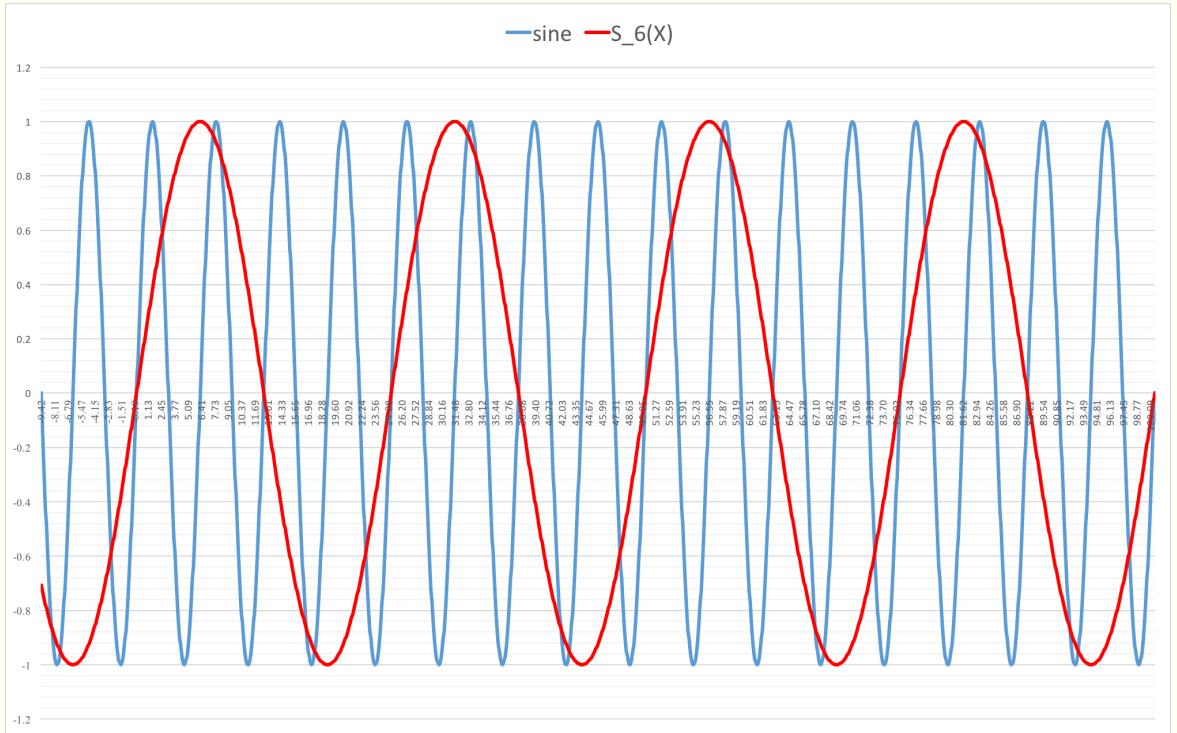
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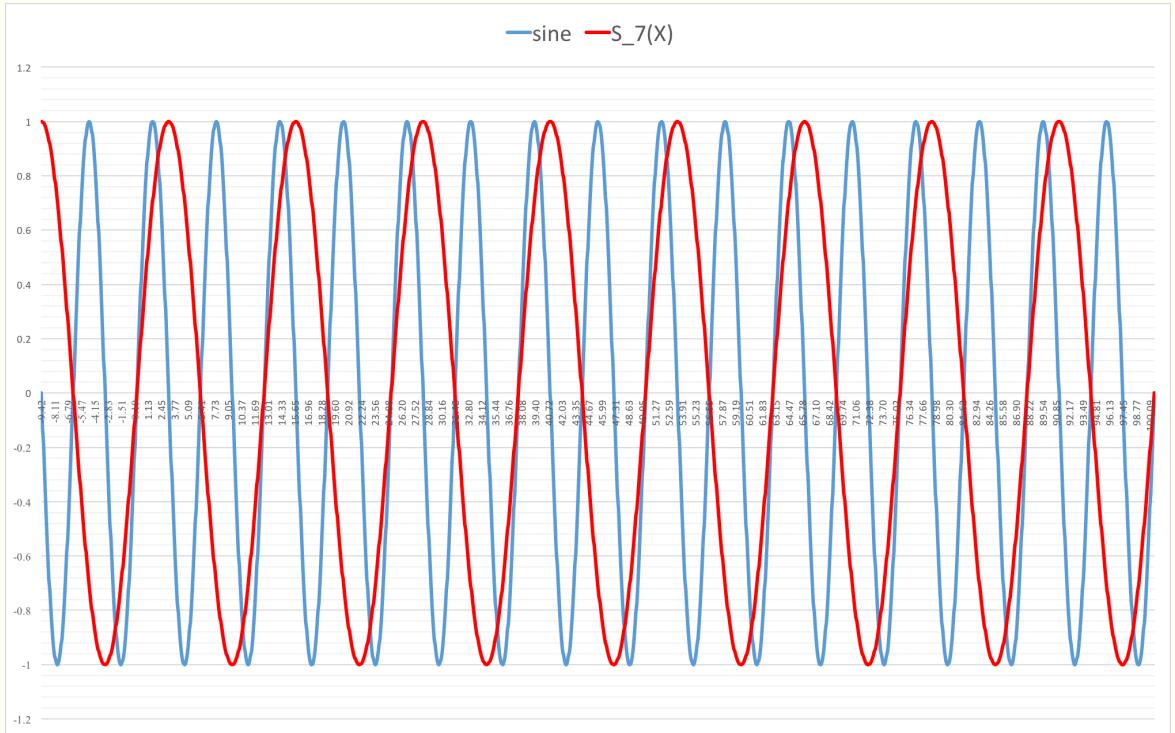
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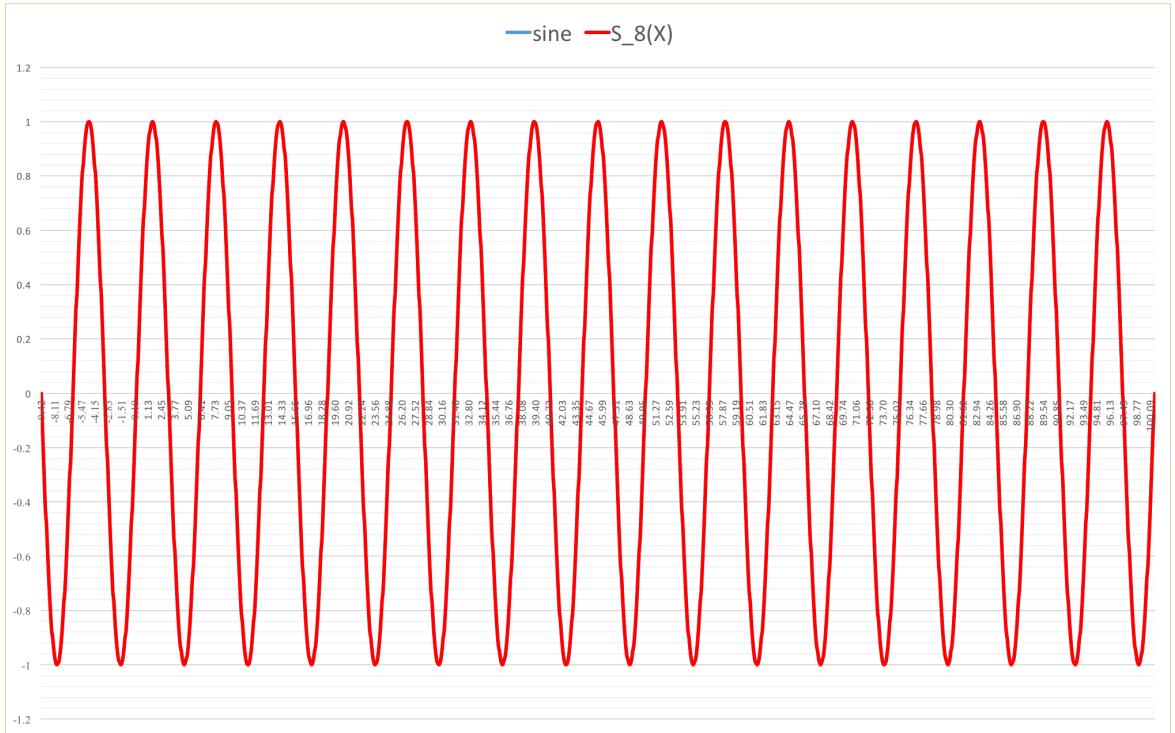
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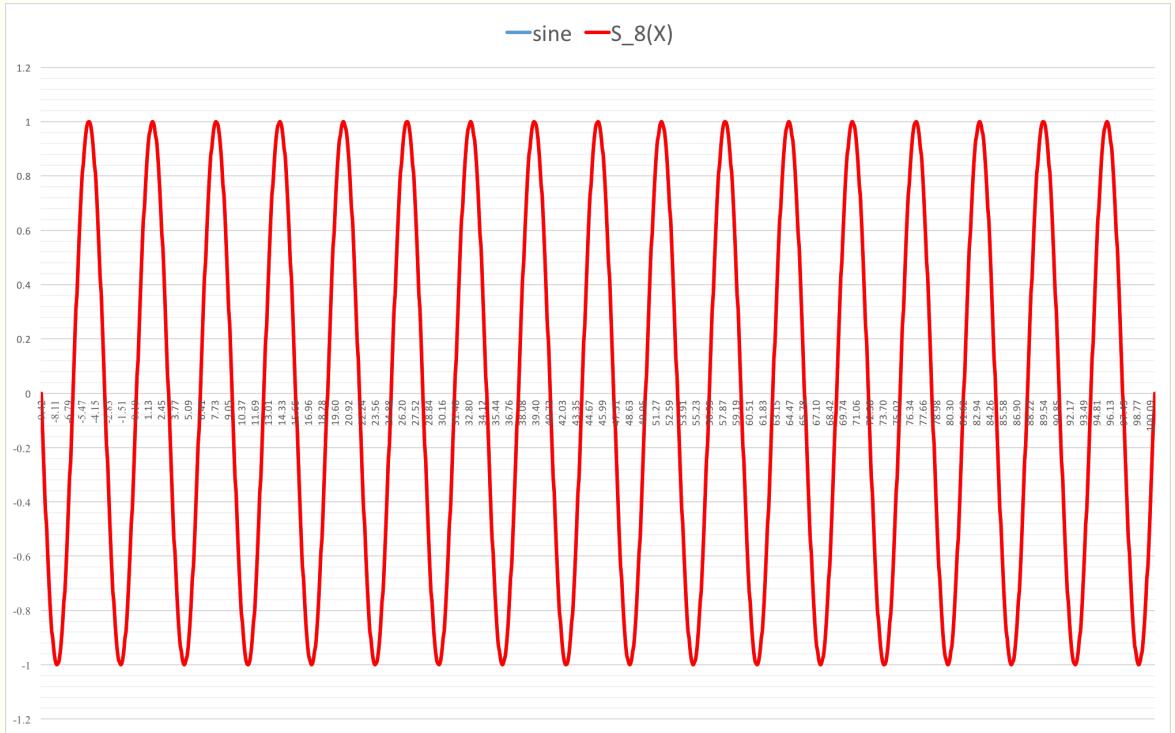
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- Numerically stable & Linear complexity

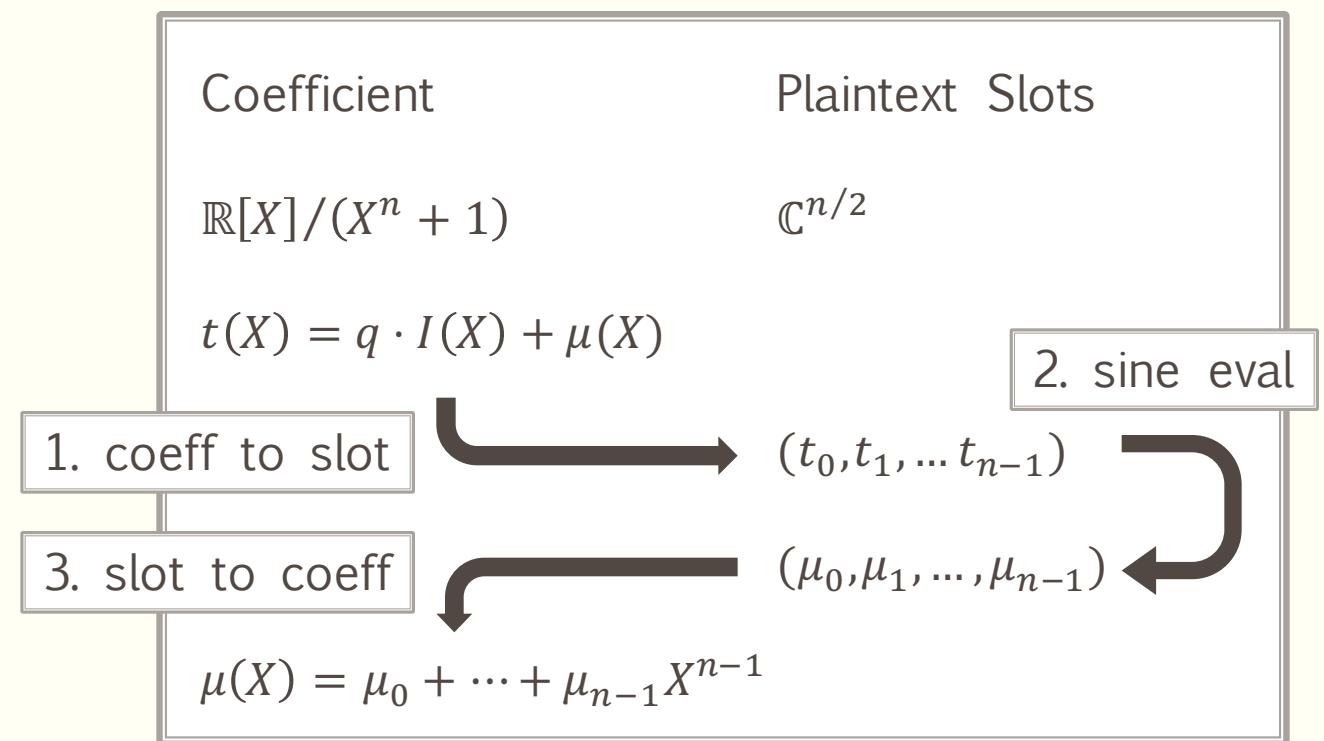


Slot-Coefficient Switching

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 - Homomorphic operations on plaintext slots, not on coefficients
 - We need to perform the modulo reduction on coefficients

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- Performance of Bootstrapping
 - Depth consumption : Sine evaluation
 - Complexity: Slot-Coefficient switchings (# of slots)

- Experimental Results
 - $127 + 12 = 139$ s / 128 slots X 12 bits
 - $456 + 68 = 524$ s / 128 slots X 24 bits

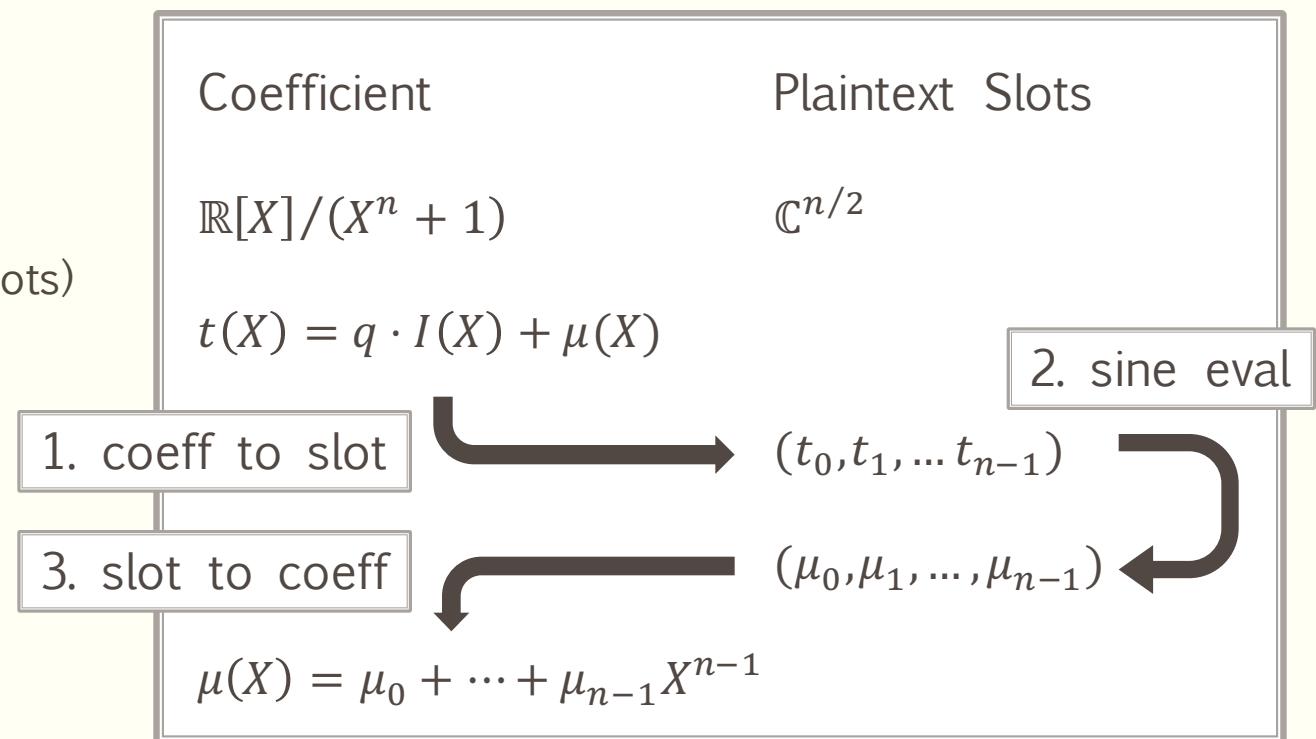


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- **Background**
- **Construction**
 - [CKKS, AC17] Homomorphic Encryption for Arithmetic of Approximate Numbers
- **Bootstrapping**
 - [CHKKS, EC18] Bootstrapping for Approximate Homomorphic Encryption
- **Related Works**

Followed Work

- Improved Bootstrapping for Approximate Homomorphic Encryption
 - Joint work with Hao Chen and Ilaria Chillotti (submission to EC19)
 - FFT-like algorithms to optimize Slot-Coefficient switchings
 - Better evaluation of sine function based on Chebyshev approximation
- [JKLS, CCS18] Secure Outsourced Matrix Computation and Application to Neural Networks
 - Evaluation of an encrypted CNN model on the encrypted MNIST data
- [CHKKS, SAC18] A Full RNS Variant of Approximate Homomorphic Encryption
 - Better performance without any high-precision arithmetic library
 - iDASH 2018
- [KS, ICISC18] Approximate Homomorphic Encryption over the Real Numbers

